

Networks and Processes

Sample Solutions to Exercise 7

Discussion on – 22.1.2009, 15h45 (submissions of solutions are highly recommended, either before or after the exercise session)

1.
 - a) Equivalent.
 - b) Not equivalent.
 - c) Not equivalent.
 - d) Equivalent.

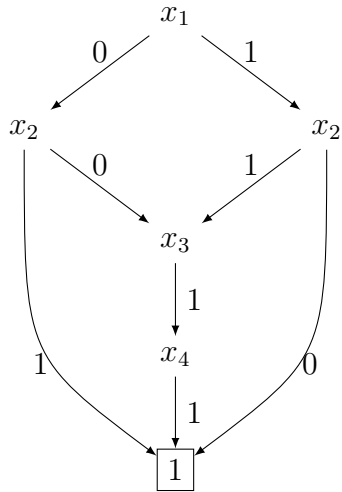
2.
 - (i) Proof of (1).
 If a Kripke structure $\mathcal{K} \models \mathbf{EF AG} \phi$,
 \Rightarrow there exists a run ρ starting at s^0 , $\exists i \geq 0$ s.t. $\rho(i) \in \llbracket \mathbf{AG} \phi \rrbracket$.
 $\Rightarrow \forall j \geq i : \rho(j) \in \llbracket \phi \rrbracket$.
 $\Rightarrow \rho(i) \in \llbracket \phi \rrbracket$.
 $\Rightarrow \mathcal{K} \models \mathbf{EF} \phi$.
 - (ii) Proof of (2).
 If a Kripke structure $\mathcal{K} \models \mathbf{AG} \phi$,
 \Rightarrow for all runs ρ starting at s^0 , $\forall i \geq 0$ s.t. $\rho(i) \in \llbracket \phi \rrbracket$.
 $\Rightarrow \rho(i) \in \llbracket \mathbf{AG} \phi \rrbracket$.
 $\Rightarrow \rho(i) \in \llbracket \mathbf{EF AG} \phi \rrbracket$.
 $\Rightarrow \mathcal{K} \models \mathbf{AG EF AG} \phi$.
 - (iii) $\mathbf{EF AG EF AG} \phi \stackrel{(1)}{\Rightarrow} \mathbf{EF EF AG} \phi \equiv \mathbf{EF AG} \phi$.
 - (iv) $\mathbf{AG} \phi \stackrel{(2)}{\Rightarrow} \mathbf{AG EF AG} \phi$, so $\mathbf{EF AG} \phi \Rightarrow \mathbf{EF AG EF AG} \phi$.
 - (v) From (iii) and (iv), $\mathbf{EF AG EF AG} \phi \equiv \mathbf{EF AG} \phi$.
 - (vi) From (v), since ϕ is an arbitrary CTL formula, we can negate both sides of the equivalence and obtain $\mathbf{AG EF AG EF} \phi \equiv \mathbf{AG EF} \phi$.

3.
 - a) $\llbracket p \mathbf{EU} q \rrbracket = S \setminus \{s_6\}$.
 - b) $\llbracket r \wedge (p \mathbf{AU} q) \rrbracket = \{s_2, s_5, s_8\}$.
 - c) $\llbracket \mathbf{EX} (r \wedge (p \mathbf{AU} q)) \rrbracket = \{s_1, s_2, s_5, s_7, s_8\}$.
 - d) $\llbracket \mathbf{AF} q \rrbracket = \{s_1, s_2, s_4, s_5, s_8\}$.
 - e) $\llbracket \mathbf{EG}_f (q \vee r) \rrbracket = \{s_0, s_4, s_6, s_7\}$. Refer to the slide page 415.
 - (i) $\mathcal{K}_{q \vee r}$ contains every state of \mathcal{K} except s_3 .
 - (ii) The non-trivial SCCs are $\{s_0, s_6\}, \{s_1, s_2\}, \{s_4\}, \{s_5\}, \{s_7\}$.

(iii) The non-trivial SCCs that contain a state from the fairness constrain are $\{s_0, s_6\}$, $\{s_4\}$, and $\{s_7\}$.

(iv) $\mathbf{EG}_f(q \vee r)$ contains states in $\mathcal{K}_{q \vee r}$ that can reach $\{s_0, s_4, s_6, s_7\}$.

4. a)



b)

